

A Brief Introduction to Using LLVM

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 - Machine code generation libraries
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 - A simple, typed IR (*bitcode*)
 - Program analysis / optimization libraries
 - Machine code generation libraries
 - Tools that compose the libraries to perform tasks
- Easy to add / remove / change functionality

What is LLVM Bitcode?

- A (Relatively) Simple IR

```
#include<stdio.h>

void
foo(unsigned e) {
    for (unsigned i = 0; i < e; ++i) {
        printf("Hello\n");
    }
}

int
main(int argc, char **argv) {
    foo(argc);
    return 0;
}
```

Code

clang -c -S -emit-llvm -O1 -g0

@str = private constant [6 x i8] c"Hello\00"

define void @foo(i32) {
 %2 = icmp eq i32 %0, 0
 br il %2, label %3, label %4

; <label>:3: ; preds = %4, %1
 ret void

; <label>:4: ; preds = %1, %4
 %5 = phi i32 [%7, %4, [0, %1]]
 %6 = tail call i32 @preamble(i8* getelementptr
 ([6 x i8], [6 x i8]* @str, i64 0, i64 0))
 %7 = add nuw i32 %5, 1
 %8 = icmp eq i32 %7, %0
 br il %8, label %3, label %4
}

define i32 @main(i32, i8** nocapture readnone) {
 tail call void @foo(i32 %0)
 ret i32 0
}

IR

What is LLVM Bitcode?

- A (Relatively) Simple IR

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#include<stdio.h>

void
foo(unsigned e) {
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main(int argc, char **argv) {
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clang -c -emit-llvm
(and llvm-dis)

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    %6 = tail call i32 @puts(i8* getelementptr
        ([6 x i8], [6 x i8]* @str, i64 0, i64 0))
    %7 = add nuw i32 %5, 1
    %8 = icmp eq i32 %7, %0
    br il %8, label %3, label %4

}

define i32 @main(i32, i8** nocapture readnone) {
    tail call void @foo(i32 %0)
    ret i32 0
}
```

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#include<stdio.h>
```

```
void
foo(unsigned e) {
    for (unsigned i = 0; i < e; ++i) {
        printf("Hello\n");
    }
}
```

```
int
main(int argc, char **argv) {
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Functions

```
@str = private constant [6 x i8] c"Hello\00"

define void @foo(i32) {
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    br il %2, label %3, label %4

; <label>:3:                                     ; preds = %4, %1
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; <label>:4:                                     ; preds = %1, %4
    %5 = phi i32 [ %7, %4 ], [ 0, %1 ]
    %6 = tail call i32 @puts(i8* getelementptr
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    %7 = add nuw i32 %5, 1
    %8 = icmp eq i32 %7, %0
    br il %8, label %3, label %4
}

define i32 @main(i32, i8** nocapture readonly) {
    tail call void @foo(i32 %0)
    ret i32 0
}
```

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```

Basic Blocks

```
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define i32 @main(i32, i8** nocapture readnone) {
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```

Basic Blocks

labels & predecessors

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; <label>:3:                                     ; preds = %4, %1
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; <label>:4:                                     ; preds = %1, %4
    %5 = pni i32 [ %7, %4 ], [ 0, %1 ]
    call i32 @puts(i8* getelementptr
                    i8, [6 x i8]* @str, i64 0, i64 0))
    %7 = add nuw i32 %5, 1
    %8 = icmp eq i32 %7, %0
    br il %8, label %3, label %4

define i32 @main(i32, i8** nocapture readonly) {
    tail call void @foo(i32 %0)
    ret i32 0
}
```

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foo(unsigned e) {
    for (unsigned i = 0; i < e; ++i) {
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    }
}

int
main(int argc, char **argv) {
    foo(argc);
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}
```

Basic Blocks

branches & successors

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@str = private constant [6 x i8] c"Hello\00"

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; <label>:3:                                ; preds = %4, %1
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; preds = %1, %4
    %5 = icmp i32 [ %7, %4 ], [ 0, %1 ]
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Instructions

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    %7 = add nuw i32 %5, 1
    %8 = icmp eq i32 %7, %0
    br i1 %8, label %3, label %4

}

define i32 @main(i32, i8** nocapture readonly) {
    tail call void @foo(i32 %0)
    ret i32 0
}
```

Inspecting Bitcode

- LLVM libraries help examine the bitcode
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```
Module& module = ...;  
for (Function& fun : module) {  
    for (BasicBlock& bb : fun) {  
        for (Instruction& i : bb) {  
  
...  
    }  
}
```

Iterate over the:

- Functions in a Module
- BasicBlocks in a Function
- Instructions in a BasicBlock

Inspecting Bitcode

- LLVM libraries help examine the bitcode
 - Easy to examine and/or manipulate
 - Many helpers (e.g. CallSite,)

```
Module& module = ...;
for (Function& fun : module) {
    for (BasicBlock& bb : fun) {
        for (Instruction& i : bb) {
            CallSite cs(&i);
            if (!cs.getInstruction()) {
                continue;
            }
        }
    }
}
```

CallSite helps you extract information from Call and Invoke instructions.

...

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- LLVM libraries help examine the bitcode
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Module &module = ...;
for (Function& fun : module) {
    for (BasicBlock& bb : fun) {
        for (Instruction& i : bb) {
            CallSite cs(&i);
            if (!cs.getInstruction()) {
                continue;
            }
            outs() << "Found a function call: " << i << "\n";
        ...
    }
}
```

Inspecting Bitcode

- LLVM libraries help examine the bitcode
 - Easy to examine and/or manipulate
 - Many helpers (e.g. CallSite, outs(), dyn_cast)

```
Module &module = ...;
for (Function& fun : module) {
    for (BasicBlock& bb : fun) {
        for (Instruction& i : bb) {
            CallSite cs(&i);
            if (!cs.getInstruction()) {
                continue;
            }
            outs() << "Found a function call: " << i << "\n";
            Value* called = cs.getCalledValue()->stripPointerCasts();
            if (Function* f = dyn_cast<Function>(called)) {
                outs() << "Direct call to function: " << f->getName() << "\n";
            }
        }
    }
}
```

dyn_cast() efficiently checks the runtime types of LLVM IR components.

Dealing with SSA

- You may ask where certain values came from
 - Useful for tracking dependencies (PDG)
 - “Where was this variable defined?”

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```
void foo()
  unsigned i = 0;
  while (i < 10) {
    i = i + 1;
  }
}
```

Dealing with SSA

- You may ask where certain values came from
- LLVM IR provides this through SSA form

```
void foo()
  unsigned i = 0;
  while (i < 10) {
    i = i + 1;
  }
}
```

What is the single definition
of *i* at this point?

Dealing with SSA

- Thus the phi (φ) instruction
 - It selects which of the definitions to use
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void foo() {  
    unsigned i = 0;  
    while (i < 10) {  
        i = i + 1;  
    }  
}
```

```
define void @foo() {  
    br label %1  
  
; <label>:1 ; preds = %1, %0  
    %i.phi = phi i32 [ 0, %0 ], [ %2, %1 ]  
    %2 = add i32 %i.phi, 1  
    %exitcond = icmp eq i32 %2, 10  
    br i1 %exitcond, label %3, label %1  
  
; <label>:3 ; preds = %1  
    ret void  
}
```

Dealing with SSA

- Thus the phi (φ) instruction
 - It selects which of the definitions to use
 - Always at the start of a basic block

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void foo()
  unsigned i = 0;
  while (i < 10) {
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}
```

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  br label %1

; <label>:1 ; preds = %1, %0
  %i.phi = phi i32 [ 0, %0 ], [ %2, %1 ]
  %2 = add i32 %i.phi, 1
  %exitcond = icmp eq i32 %2, 10
  br i1 %exitcond, label %3, label %1

; <label>:3 ; preds = %1
  ret void
}
```

Dependencies in General

- You can loop over the values an instruction uses

```
for (Use& u : inst->operands()) {  
    // inst uses the Value* u  
}
```

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```
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}
```

for %a = %b + %c:

[%b, %c]

Dependencies in General

- You can loop over the values an instruction uses

```
for (Use& u : inst->operands()) {  
    // inst uses the Value* u  
}
```

- You can loop over the instructions that use a particular value

```
Instruction* inst = ...;  
for (User* user : inst->users())  
    if (auto* i = dyn_cast<Instruction>(user)) {  
        // inst is used by Instruction i  
    }
```

Dealing with Types

- LLVM IR is *strongly typed*
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```
define i64 @trunc(i16 zeroext %a) {  
    %1 = zext i16 %a to i64  
    ret i64 %1  
}
```

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define i64 @trunc(i16 zeroext %a) {  
    %1 = zext i16 %a to i64  
    ret i64 %1  
}
```

- Also types for pointers, arrays, structs, etc.
 - Strong typing means they take a bit more work

Dealing with Types: GEP

- We sometimes need to extract elements/fields from arrays/structs
 - Pointer arithmetic
 - Done using GetElementPointer (GEP)

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```
struct rec {  
    int x;  
    int y;  
};  
  
struct rec *buf;  
  
void foo() {  
    buf[5].y = 7;  
}
```

Dealing with Types: GEP

- We sometimes need to extract elements/fields from arrays/structs
 - Pointer arithmetic
 - Done using GetElementPointer (GEP)

```
%struct.rec = type { i32, i32 }
@buf = global %struct.rec* null
```

```
define void @foo() {
    %1 = load %struct.rec*, %struct.rec** @buf
    %2 = getelementptr %struct.rec, %struct.rec* %1, i64 5, i32 1
    store i32 7, i32* %2
    ret void
}
```

```
struct rec {
    int x;
    int y;
};

struct rec *buf;

void foo() {
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- We sometimes need to extract elements/fields from arrays/structs
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define void @foo() {
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    store i32 7, i32* %2
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}
```

```
struct rec {
    int x;
    int y;
};

struct rec *buf;

void foo() {
    buf[5].y = 7;
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Where Can You Get Info?

- The online documentation is extensive:
 - LLVM Programmer's Manual
 - LLVM Language Reference Manual

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 - LLVM Programmer's Manual
 - LLVM Language Reference Manual
- The header files!
 - All in `llvm-3.x.src/include/llvm/`

`BasicBlock.h`

`CallSite.h`

`DerivedTypes.h`

`Function.h`

`Instructions.h`

`InstrTypes.h`

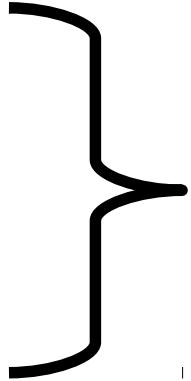
`IRBuilder.h`

`Support/InstVisitor.h`

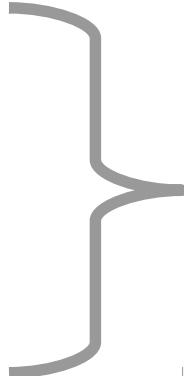
`Type.h`

Creating a *Static Analysis*

Making a New Analysis

- Analyses are organized into individual *passes*
 - ModulePass
 - FunctionPass
 - LoopPass
 - ...
- 
- Derive from the appropriate base class to make a Pass

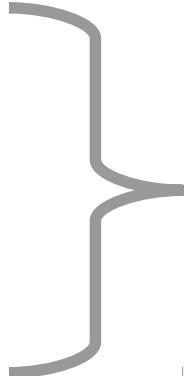
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3 Steps

- 1) Declare your pass
- 2) Register your pass
- 3) Define your pass

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3 Steps

- 1) Declare your pass
- 2) Register your pass
- 3) Define your pass

Let's count the number of **static direct calls** to each function.

Making a ModulePass (1)

- Declare your ModulePass

```
struct StaticCallCounter : public llvm::ModulePass {

    static char ID;

    DenseMap<Function*, uint64_t> counts;

    StaticCallCounter()
        : ModulePass(ID)
    { }

    bool runOnModule(Module& m) override;

    void print(raw_ostream& out, const Module* m) const override;

    void handleInstruction(CallSite cs);
};
```

Making a ModulePass (1)

- Declare your ModulePass

```
struct StaticCallCounter : public llvm::ModulePass {  
  
    static char ID;  
  
    DenseMap<Function*, uint64_t> counts;  
  
    StaticCallCounter()  
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    bool runOnModule(Module& m) override;  
  
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};
```

Making a ModulePass (1)

- Declare your ModulePass

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struct StaticCallCounter : public llvm::ModulePass {  
  
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    bool runOnModule(Module& m) override;  
  
    void print(raw_ostream& out, const Module* m) const override;  
  
    void handleInstruction(CallSite cs);  
};
```

Making a ModulePass (2)

- Register your ModulePass
 - This allows it to be dynamically loaded as a plugin

```
char StaticCallCounter::ID = 0;  
  
RegisterPass<StaticCallCounter> SCCReg("callcounter",  
                                         "Print the static count of direct calls");
```

Making a ModulePass (3)

- Define your ModulePass
 - Need to override `runOnModule()` and `print()`

```
bool  
StaticCallCounter::runOnModule(Module& m) {  
    for (auto& f : m)  
        for (auto& bb : f)  
            for (auto& i : bb)  
                handleInstruction(CallSite(&i));  
    return false; // False because we didn't change the Module  
}
```

Making a ModulePass (3)

- analysis continued...

```
void
StaticCallCounter::handleInstruction(CallSite cs) {
    // Check whether the instruction is actually a call
    if (!cs.getInstruction()) { return; }

    // Check whether the called function is directly invoked
    auto called = cs.getCalledValue()->stripPointerCasts();
    auto fun    = dyn_cast<Function>(called);
    if (!fun) { return; }

    // Update the count for the particular call
    auto count = counts.find(fun);
    if (counts.end() == count) {
        count = counts.insert(std::make_pair(fun, 0)).first;
    }
    ++count->second;
}
```

Making a ModulePass (3)

- analysis continued...

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- analysis continued...

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    }
    ++count->second;
}
```

Making a ModulePass (3)

- Printing out the results

```
void
CallCounterPass::print(raw_ostream& out, const Module* m) const {
    out << "Function Counts\n"
        << "=====\\n";
    for (auto& kvPair : counts) {
        auto* function = kvPair.first;
        uint64_t count = kvPair.second;
        out << function->getName() << " : " << count << "\\n";
    }
}
```

Creating a *Dynamic* Analysis

Making a Dynamic Analysis

- We've counted the static direct calls to each function.
- How might we compute the ***dynamic calls*** to each function?

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- Need to *modify* the original program!

Making a Dynamic Analysis

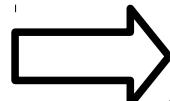
- We've counted the static direct calls to each function.
- How might we compute the *dynamic calls* to each function?
- Need to *modify* the original program!
- Steps:
 - 1) **Modify** the program using passes
 - 2) **Compile** the modified version
 - 3) **Run** the new program

Modifying the Original Program

Goal: Count the dynamic calls to each function in an execution.

So how do we want to modify the program?

```
void foo()  
    bar();  
}
```

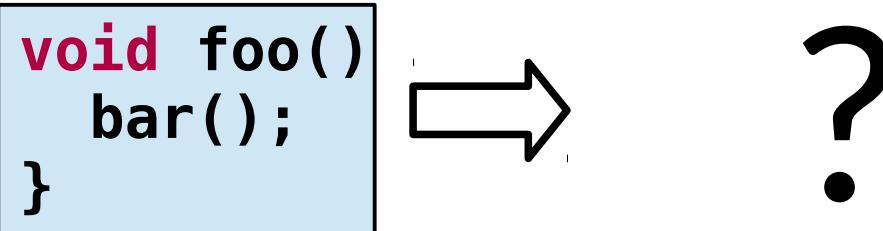


?
•

Modifying the Original Program

Goal: Count the dynamic calls to each function in an execution.

So how do we want to modify the program?

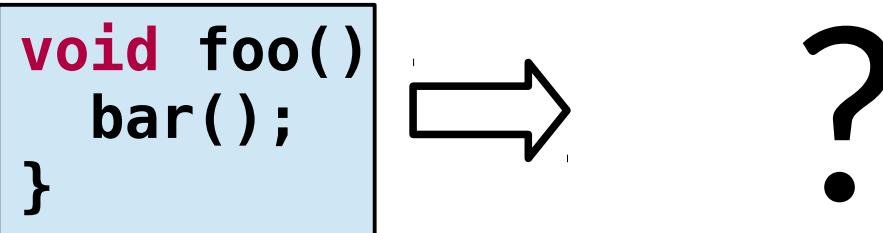


Keep a counter for each function!

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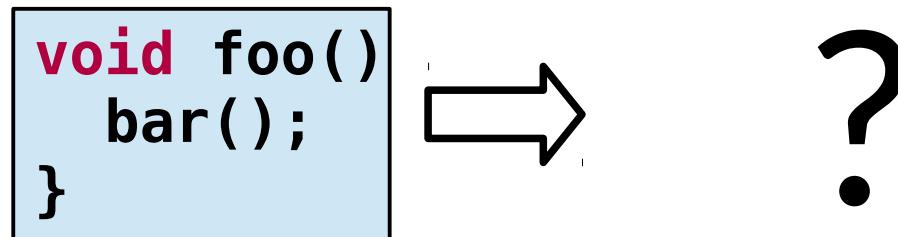
Keep a counter for each function!

2 Choices:

Modifying the Original Program

Goal: Count the dynamic calls to each function in an execution.

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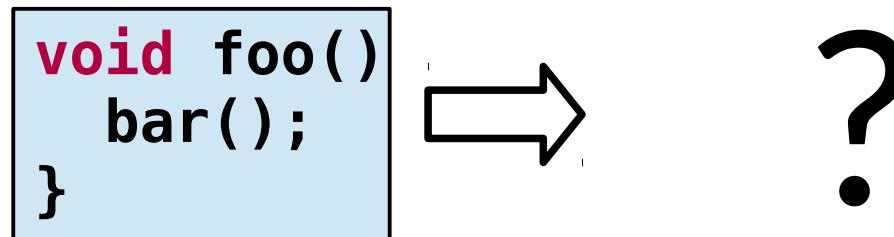
2 Choices:

- 1) increment count for each function *as it starts*
- 2) increment count for each function *at its call site*

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2 Choices:

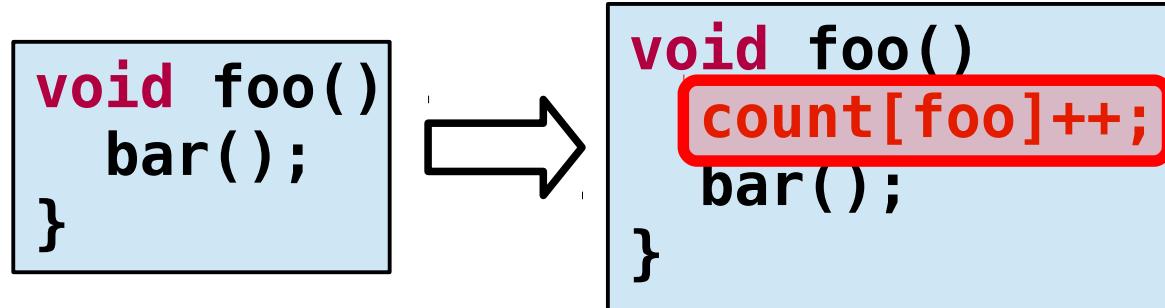
- 1) increment count for each function *as it starts*
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Does that even matter? Are there trade offs?

Modifying the Original Program

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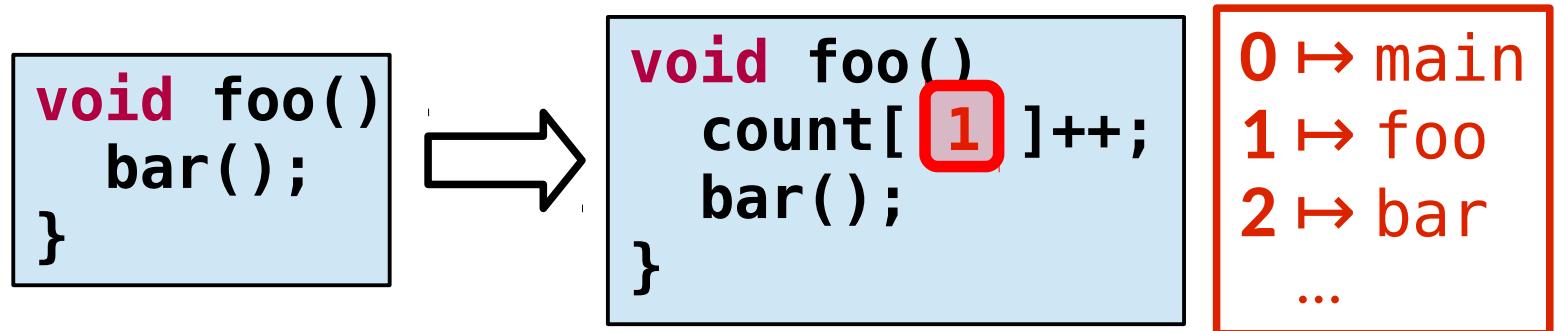


We'll increment at the function entry.
(The demo code has both)

Modifying the Original Program

Goal: Count the dynamic calls to each function in an execution.

So how do we want to modify the program?



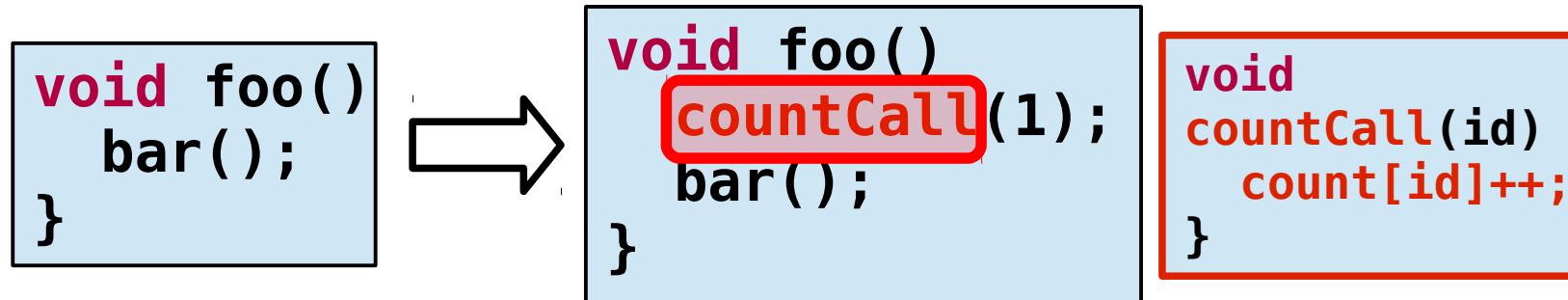
We'll increment at the function entry

- *Using numeric IDs* for functions is sometimes easier

Modifying the Original Program

Goal: Count the dynamic calls to each function in an execution.

So how do we want to modify the program?



We'll increment at the function entry

- Using numeric IDs for functions is sometimes easier
- Inserting function calls is sometimes easier

Modifying the Original Program

What might adding this call look like?

```
void DynamicCallCounter::handleInstruction(CallSite cs, Value* counter) {
    // Check whether the instruction is actually a call
    if (!cs.getInstruction()) {
        return;
    }

    // Check whether the called function is directly invoked
    auto calledValue = cs.getCalledValue()->stripPointerCasts();
    auto calledFunction = dyn_cast<Function>(calledValue);
    if (!calledFunction) {
        return;
    }

    // Insert a call to the counting function.
    IRBuilder<> builder(cs.getInstruction());
    builder.CreateCall(counter, builder.getInt64(ids[calledFunction]));
}
```

Modifying the Original Program

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    // Check whether the instruction is actually a call
    if (!cs.getInstruction()) {
        return;
    }

    // Check whether the called function is directly invoked
    auto intercasts = cs.getInstruction()->getInterCasts();
    auto callees = cs.getInstruction()->getFunction()->getFunctionCallees();
    if (!callees || !callees->empty()) {
        return;
    }

    // Insert a call to the counting function.
    IRBuilder<> builder(cs.getInstruction());
    builder.CreateCall(counter, builder.getInt64(ids[calledFunction]));
}
```

In practice, it's more complex.

You can find details in the demo code.

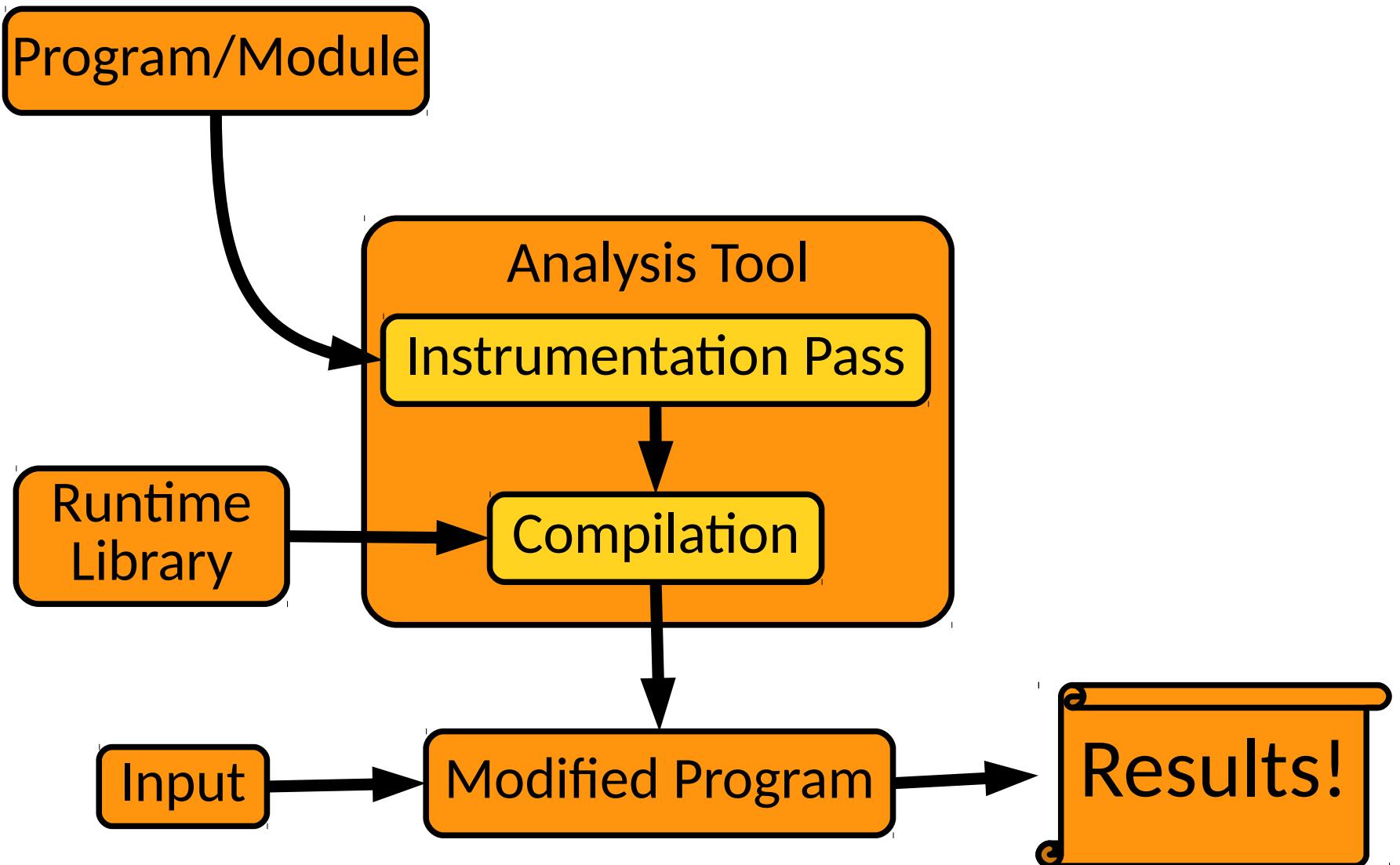
Using a Runtime Library

Don't forget that we need to put countCall()
somewhere!

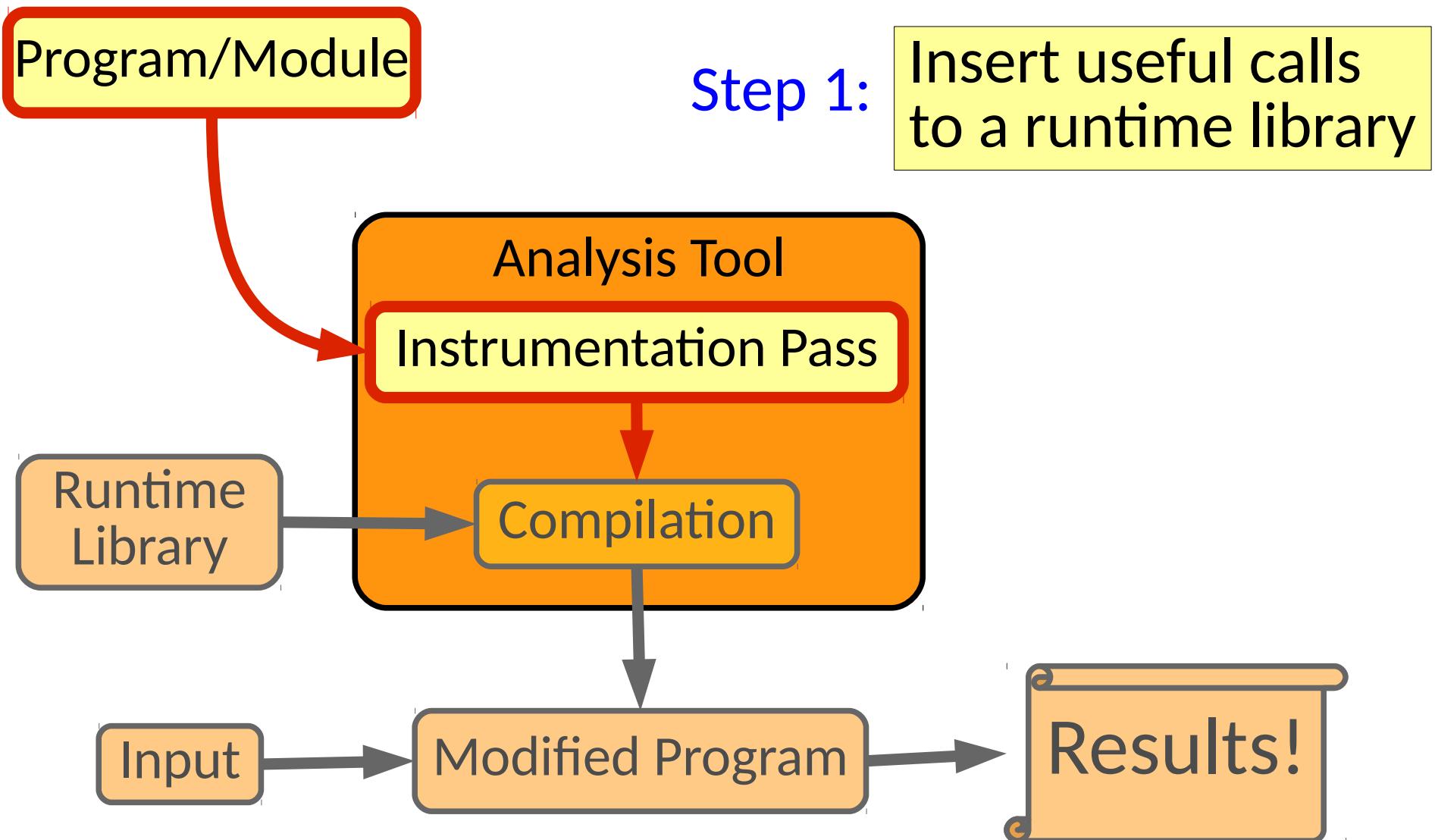
- Placed in a library linked with the main executable

```
void  
countCalled(uint64_t id) {  
    ++functionInfo[id];  
}
```

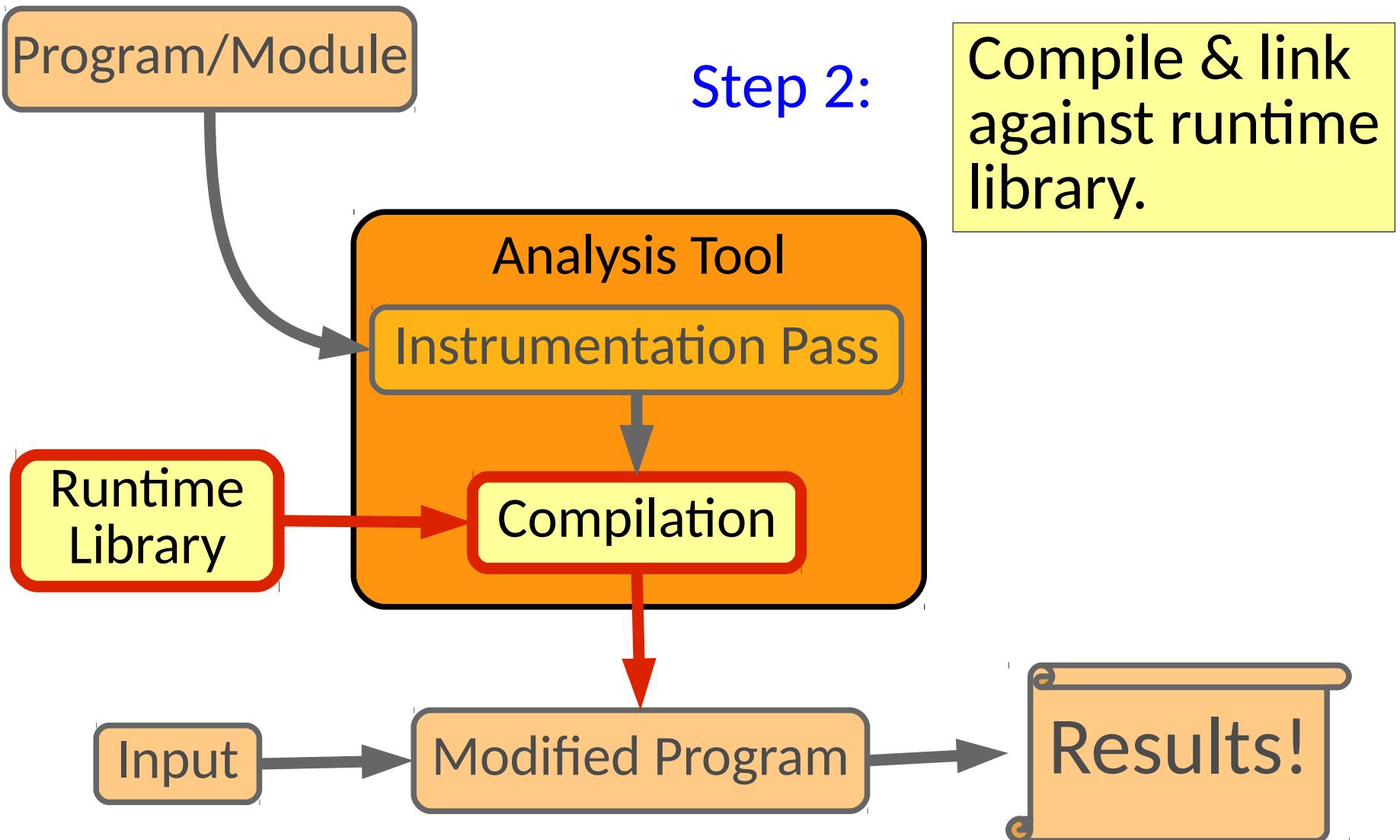
Dynamic Analysis Big Picture



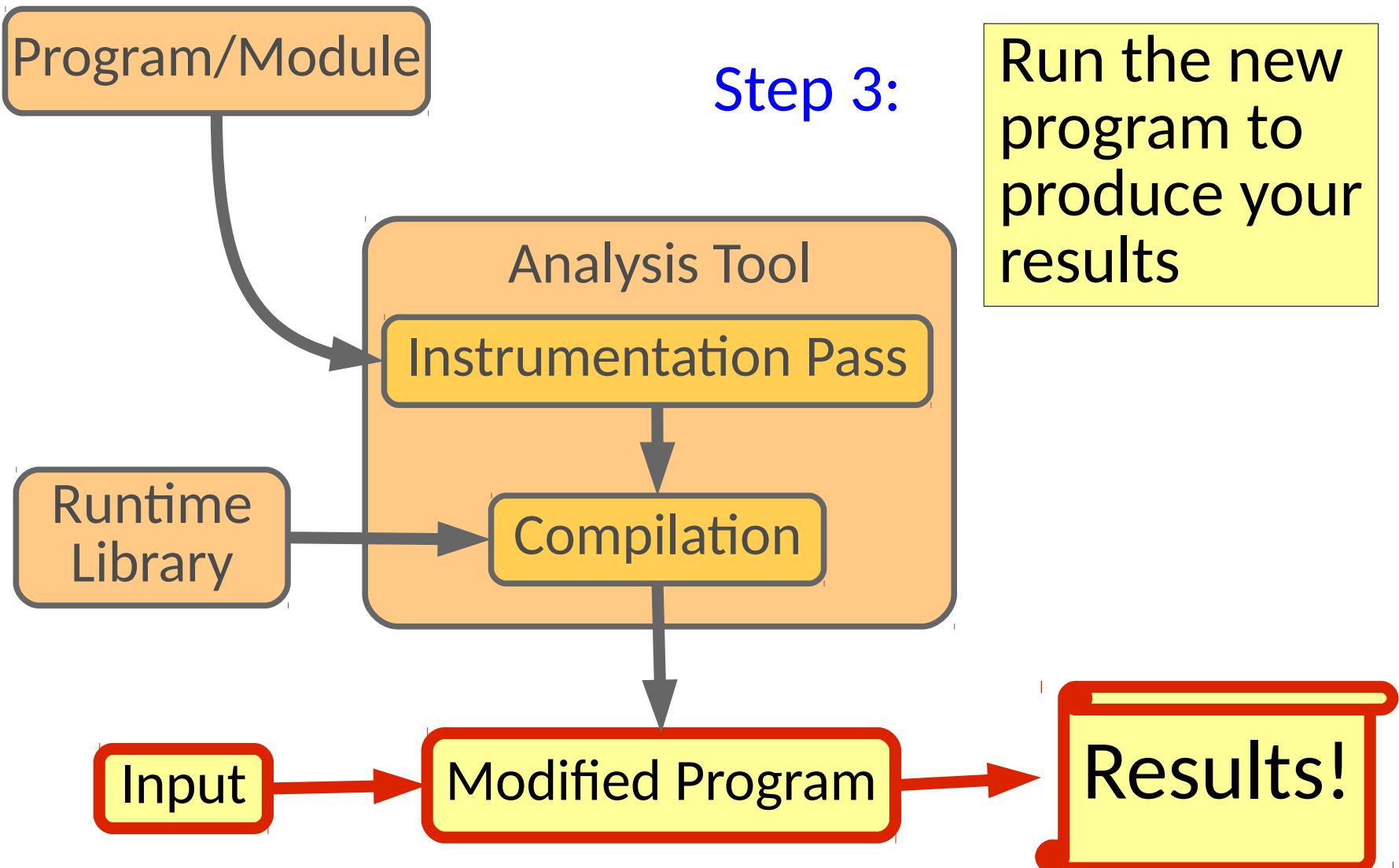
Dynamic Analysis Big Picture



Dynamic Analysis Big Picture



Dynamic Analysis Big Picture



Bringing It All Together

LLVM Projects

- LLVM organizes groups of passes and tools into *projects*

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- LLVM organizes groups of passes and tools into *projects*
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- For the most part, you can follow the directions online & in project description

Extra Tips

- I have a pointer to something. What is it?
 - The getName() method works on most things.
 - You can usually: outs() << x

Extra Tips

- I have a pointer to something. What is it?
 - The getName() method works on most things.
 - You can usually: outs() << x
- Sadly no longer true:

How do I see the C++ API calls for constructing a module?

- `llc -march=cpp <bitcode>.bc -o <cppapi>.cpp`